

Specification of state-based Systems

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States and Transition Systems

Goals

- Understanding the difference between explicitly and implicitly given states
- Learn to understand states and transition systems

Functional View:

- A state is represented by a constructor term
- States are passed as explicit parameters of functions

Example (Stack):

The signature of stacks over natural numbers:

sig STACKSIG=

sorts Nat, Stack
ops empty : Stack
 push : Nat × Stack → Stack
 top : Stack → Nat
 pop : Stack → Stack

end

A state z_1 is given by following constructor term:

$$z_1 =_{def} \text{push}(2, \text{push}(1, \text{empty}))$$

The invocation $\text{pop}(z_1)$ changes the **state** z_1 into the **state** $z_2 = \text{push}(1, \text{empty})$.

State-based view:

- A state is implicit and
- The operations have following types.

ops empty : Unit
push : Nat \rightarrow Unit
top : Unit \rightarrow Nat
pop : Unit \rightarrow Unit

The invocation of pop() changes the state z_1 into the state z_2 . A following push(4) leads to the state

$$z_3 =_{def} \text{push}(4, \text{push}(1, \text{empty}))$$

Example (State-based view of Person):

The description of persons requires several attributes:

name : Name

address : Address

age : Nat

i.e. a state is given by several components.

Operations:

getName : Unit \rightarrow Name

getAddress : Unit \rightarrow Address

getAge : Unit \rightarrow Nat

Example (Functional (algebraic) view of Person):

generated type $\text{Person} ::= \text{makePerson}(\text{getName: Name};$
 $\text{getAddress: Address}; \text{getAge: Nat})$

This defines the set of states with a constructor

$\text{makePerson} : \text{Name} \times \text{Address} \times \text{Nat} \rightarrow \text{Person}$

and the selection operations

$\text{getName} : \text{Person} \rightarrow \text{Name}$

$\text{getAddress} : \text{Person} \rightarrow \text{Address}$

$\text{getAge} : \text{Person} \rightarrow \text{Nat}$

Definition (State Signature):

Given a signature $\Sigma = (S, F, P)$ and a Σ -structure A .

- A **state signature** is a family of (system-) variables $(X_s)_{s \in S}$ with X_s countable for all $s \in S$.
- An **A-state** is an assignment $\sigma : X_s \rightarrow A_s$ for all $s \in S$.

Example (Person):

A state signature of Person is:

name : Name

address : Address

age : Nat

A structure A with the sorts Name, Address, Nat is e.g.

$$\begin{aligned}A_{\text{Name}} &=_{\text{def}} \text{char}^* \\A_{\text{Address}} &=_{\text{def}} \text{char}^* \times \mathbb{N} \times \text{char}^* \\A_{\text{Nat}} &=_{\text{def}} \mathbb{N}\end{aligned}$$

A state s is e.g.

$$\begin{aligned}s(\text{name}) &= \text{wirsing} \\s(\text{address}) &= (\text{Schlagintweitstr.}, 18, \text{München}) \\s(\text{age}) &= 53\end{aligned}$$

A transition system describes a set of transitions:

Definition:

Given a signature $\Sigma = (S, F, P)$ and a state signature V .

1. A **transition system** $\Gamma = (Z, \delta)$ is given by
 - a set Z of Σ -states and
 - a transition relation $\delta \subseteq Z \times Z$.
2. A **labelled transition system** $\Gamma = (Z, \mathcal{A}, \delta)$ is given by
 - a set Z of states,
 - a set \mathcal{A} of actions and
 - a transition relation $\delta \subseteq Z \times \mathcal{A} \times Z$.

Example (Stack):

Let K be a stack with

$$K_{\text{Stack}} =_{\text{def}} \mathbb{N}^*$$

and a state signature

stack : Stack

Then the set of states equals K_{Stack} , i.e.

$$Z =_{\text{def}} K_{\text{Stack}} = \mathbb{N}^*$$

the set of actions equals the invocations of stack operations, i.e.

$$A =_{\text{def}} \{\text{push}(n) \mid n \in \mathbb{N}\} \cup \{\text{pop}(), \text{top}()\}$$

and the labelled transition system has the transition relation δ with

$$(\langle \text{stack} = s \rangle, \text{push}(n), \langle \text{stack} = \langle n, s \rangle \rangle) \in \delta$$

$$(\langle \text{stack} = \langle n, s \rangle \rangle, \text{pop}(), \langle \text{stack} = s \rangle) \in \delta$$

$$(\langle \text{stack} = \langle n, s \rangle \rangle, \text{top}(), \langle \text{stack} = \langle n, s \rangle \rangle) \in \delta$$

for all $n \in \mathbb{N}$.

top has not been considered adequately in this example.

Better: Include arguments and results of the operation into the action.

$$\begin{aligned} \mathcal{A} =_{def} & \quad \mathbb{N} \times \text{push} \times \text{Unit} \\ & \cup \quad \text{Unit} \times \text{pop} \times \text{Unit} \\ & \cup \quad \text{Unit} \times \text{top} \times \mathbb{N} \end{aligned}$$

The transition relation has the form:

$$\begin{aligned} & (\langle \text{stack} = s \rangle, (n, \text{push}, ()), \langle \text{stack} = \langle n, s \rangle \rangle) \\ & (\langle \text{stack} = \langle n, s \rangle \rangle, ((), \text{pop}, ()), \langle \text{stack} = s \rangle) \\ & (\langle \text{stack} = \langle n, s \rangle \rangle, ((), \text{top}, n), \langle \text{stack} = \langle n, s \rangle \rangle) \end{aligned}$$

for all $n \in \mathbb{N}$.