Performance Modelling of Computer Systems

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Discrete-Event Simulation

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So far, performance models have been analysed through the solution of a mathematical problem (i.e., a system of linear equations, a system of coupled differential equations, etc.).

Simulation is an alternative approach which produces sample traces of the stochastic process under study.

The generation of these traces requires access to some object which is capable of providing random numbers.

When implemented on a computer, **pseudo-random** generators approximate sequences of random numbers satisfying a desired probability distribution.

Pseudo-random Generators: Java

```
/**
```

- * Returns a <code>double</code> value with a positive sign, greater
- * than or equal to <code>0.0</code> and less than <code>1.0</code>.
- * Returned values are chosen pseudorandomly with (approximately)
- * uniform distribution from that range.
- *
- * When this method is first called, it creates a single new
- * pseudorandom-number generator, exactly as if by the expression
- * <blockquote>new java.util.Random</blockquote> This
- * new pseudorandom-number generator is used thereafter for all

```
* calls to this method and is used nowhere else.
```

*

```
This method is properly synchronized to allow correct use by
more than one thread. However, if many threads need to generate
pseudorandom numbers at a great rate, it may reduce contention
for each thread to have its own pseudorandom-number generator.
```

*

```
* @return a pseudorandom <code>double</code> greater than or equal
* to <code>0.0</code> and less than <code>1.0</code>.
* @see java.util.Random#nextDouble()
```

```
*/
```

```
public static double random() {
```

```
if (randomNumberGenerator == null) initRNG();
```

```
return randomNumberGenerator.nextDouble();
```

```
}
```

rand

Uniformly distributed pseudorandom numbers

Syntax

```
r = rand(n)
rand(m,n)
rand([m,n])
rand([m,n,p,...)
rand([m,n,p,...])
rand
rand(size(A))
r = rand(..., 'double')
r = rand(..., 'single')
```

Description

r = rand(n) returns an n-by-n matrix containing pseudorandom values drawn from the standard uniform distribution on the open interval (0,1). rand(m,n) or rand([m,n]) returns an m-by-n matrix. rand(m,n,p,...) or rand([m,n,p,...]) returns an m-by-n-by-p-by-... array. rand returns a scalar. rand(size(A)) returns an array the same size as **A**.

r = rand(..., 'double') or r = rand(..., 'single') returns an array of uniform values of the specified class.

The simulator maintains a global variable, currentTime, which gives current simulated time (usually initialised to 0).

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The simulation algorithm pops the first event off the list, advances currentTime to the event's firingTime, and changes the state of the system according to the kind of event being processed.

The processing of an event may in turn generate some other event, which will be then scheduled according to its firing time.

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For ergodic models, the algorithm terminates when a given time horizon is reached (the event list will never be empty).

Example

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In the object-oriented pseudo-code that we will use, an arrival is a class which is constructed with a parameter denoting the firing time of the event.

Algorithm

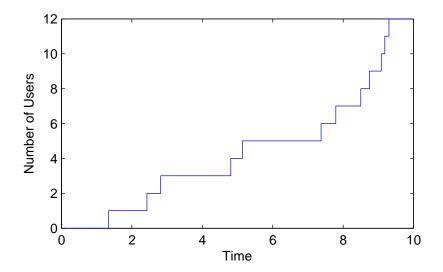
```
currentTime = 0
nUsers = 0
Trace = \{(currentTime, nUsers)\}
new Arrival(currentTime + getRandomExp())
while EventList \neq \emptyset do
  arrival = EventList.pop()
  currentTime = arrival.firingTime
  if currentTime > 10 then
     break
  end if
  nUsers = nUsers + 1
  Trace = Trace \cup {(currentTime, nUsers)}
  new Arrival(currentTime + getRandomExp())
end while
```

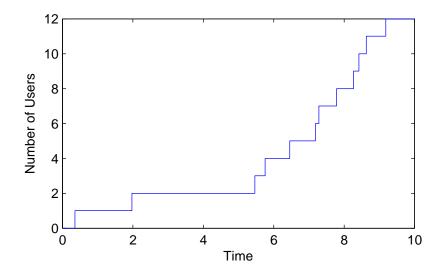
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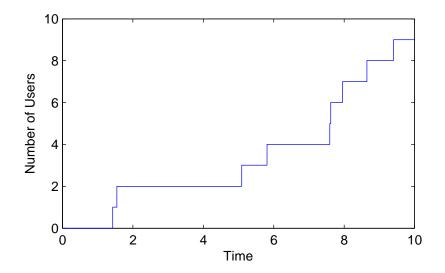
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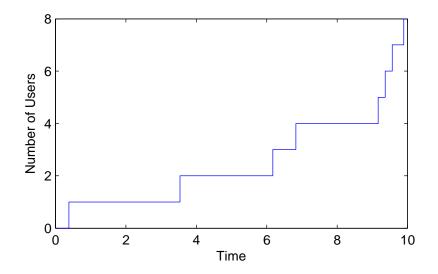
Notes:

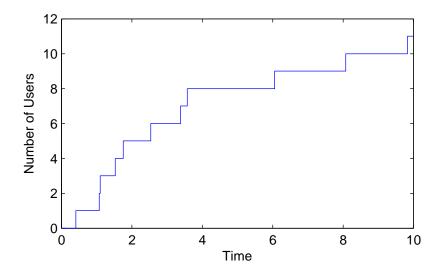
- This algorithm may be simplified
- The data structure Trace allows us to plot the sample path obtained



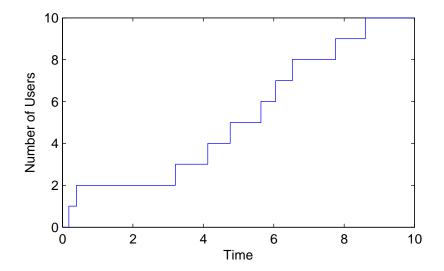


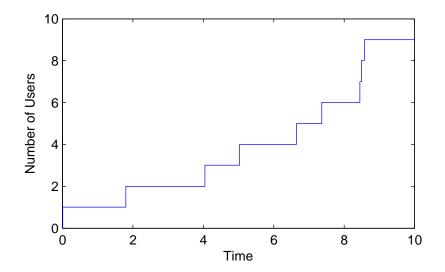


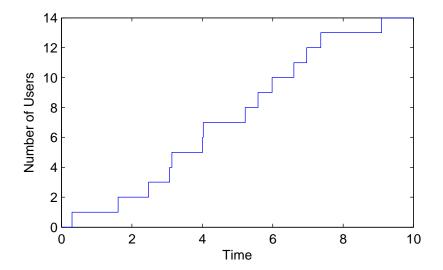


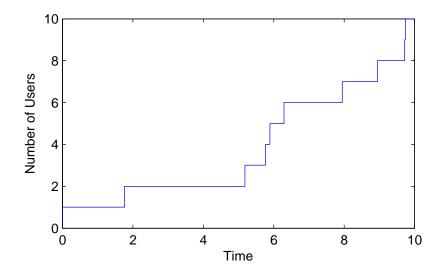


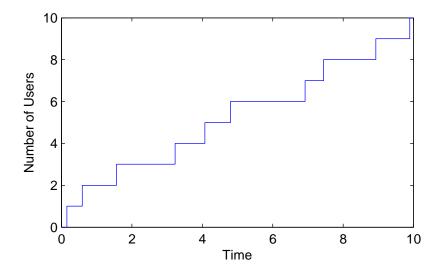
Some Traces









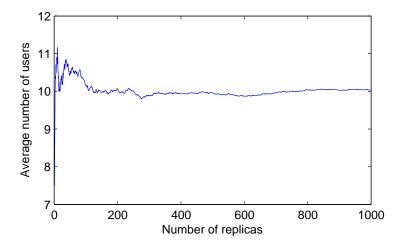


Termination Criteria

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A termination condition could be based on this error being less than some given threshold. Unfortunately, in general we do not know the exact solution — otherwise we would not simulate the model — therefore termination conditions must be based on some form of statistical analysis of the samples obtained.

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A confidence interval for the true mean m is a probability bound of the kind

$$\mathbb{P}(a < m < b) \geq \xi$$

where *a* and *b* are obtained from the data and ξ is a given confidence level (typically, $\xi = 0.95$ or $\xi = 0.99$).

Confidence Intervals

For sufficiently large *n*, an approximate confidence interval for the mean at level $1 - \alpha$ is

$$\hat{\mu} \pm \eta \frac{s_n}{\sqrt{n}},$$

where s_n is the statistical standard deviation

$$s_n^2 = \frac{1}{n} \sum_{i=1}^n (x_i - \hat{\mu}_n^2)$$

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For a desired α the corresponding value of η , η_{α} , is available in ready-to-use tables. For instance, for $\eta_{0.05} = 1.960$ (95% confidence interval) and $\eta_{0.01} = 2.576$ (99% confidence interval).

Example

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For our Poisson example, we obtained the following results:

Iteration	Mean	Error 95%	Error 99%
1000	10.190	1.882	2.474
5000	9.996	0.871	1.145
10000	9.993	0.619	9.813
15000	10.011	0.506	0.665
20000	10.004	0.438	0.575
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To halve the error requires a fourfold increase in the number of samples!

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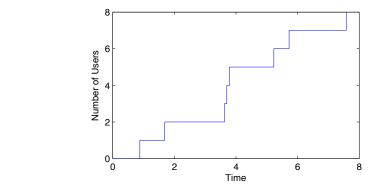
The method of batch means instead is preferred when one is interested in steady-state simulation. With the former method, the average is obtained across samples. With the latter the average is a time average.

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The simulation consists of a single long run of length T, which tracks the evolution of a performance index, I(t). The average across that index is given by

$$\hat{\mu} = rac{1}{T} \int_0^T I(t) dt$$



$$\hat{\mu} = \frac{1}{10} \times \begin{bmatrix} 0 \times 0.88 + 1 \times (1.69 - 0.88) + 2 \times (3.63 - 1.69) + \\ 3 \times (3.70 - 3.63) + 4 \times (3.79 - 3.70) + 5 \times (5.23 - 3.79) + \\ 6 \times (5.72 - 5.23) + 7 \times (7.58 - 5.72) + 7 \times (8 - 7.58) \end{bmatrix}$$

To compute confidence intervals, the single run is divided into b batches of the same length T/b and the following estimates are computed

$$\hat{\mu}_i = rac{b}{T} \int_{(i-1)T/b}^{iT/b} I(t) dt, \qquad ext{for } i = 1, \dots b$$

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Often, the model is subjected to transient removal, i.e., the statistics are not collected over some initial period of time to exclude those samples which are related to the stochastic process being away from stationary conditions.

- Stochastic Simulation for Java is a framework for discrete-event simulation for Java
- It is freely available at http://www.iro.umontreal.ca/~simardr/ssj/indexen.html
- It supports a wide array of random number generators, statistics collectors, a graph visualisation toolkit, etc.

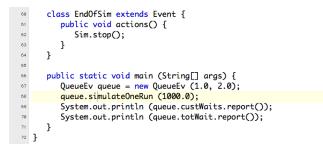
A Queue with SSJ

```
import umontreal.iro.lecuyer.simevents.*;
2 import umontreal.iro.lecuyer.rng.*;
  import umontreal.iro.lecuver.randvar.*;
 import umontreal.iro.lecuyer.stat.*;
  import java.util.LinkedList;
  public class OueueEv {
     RandomVariateGen genArr;
     RandomVariateGen aenServ:
18
     LinkedList<Customer> waitList = new LinkedList<Customer> ():
     LinkedList<Customer> servList = new LinkedList<Customer> ();
12
     Tallv custWaits = new Tally ("Waiting times");
     Accumulate totWait = new Accumulate ("Size of aueue");
14
     class Customer { double arrivTime, servTime; }
16
     public OueueEv (double lambda, double mu) {
18
        genArr = new ExponentialGen (new MRG32k3a(), lambda);
19
20
        genServ = new ExponentialGen (new MRG32k3a(), mu);
     3
22
     public void simulateOneRun (double timeHorizon) {
23
        Sim.init();
24
        new EndOfSim().schedule (timeHorizon);
25
        new Arrival().schedule (genArr.nextDouble());
26
        Sim.start();
     }
28
```

A Queue with SSJ



Tribastone (IFI LMU)



- Discrete-event simulation does not make assumptions on the simulated system.
- In particular, no assumptions are made on the probability distributions.
- Monte Carlo simulation is an algorithm for continuous-time Markov chain, which exploits the properties of the exponential distribution.

Monte Carlo Simulation

In a homogeneous CTMC,

$$q_{ij} = \lim_{\Delta t o 0} rac{p_{ij}(t,t+\Delta t)}{\Delta t}, \qquad ext{for all } t.$$

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Given a state *i*, the holding time is exponentially distributed with rate $\sum_{j} q_{ij} = -q_{ii}$. Drawing *x* from a uniform distribution in (0, 1) and computing

$$y=\frac{1}{q_{ii}}\ln(1-x)$$

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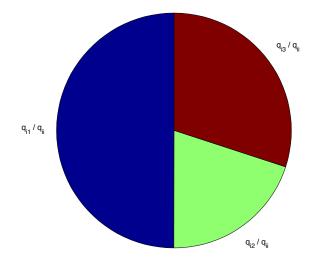
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The probability that the transition from *i* to *k* happens is given by $q_{ik} / \sum_j q_{ij}$. A sample is obtained by drawing *z* from a uniform distribution in (0, 1) and choosing the smallest *k* such that $z > \frac{\sum_{i=1}^{k-1} q_{ij}}{\sum_i q_{ij}}$.

Transition Selection



t = 0

 $s = s_1$

while termination criteria are not satisfied do

draw x, y in
$$U(0, 1)$$

 $\Delta t = \frac{1}{q_{ss}} \ln(1 - x)$
chose smallest k such that $y > \frac{\sum_{j=1}^{k-1} q_{sj}}{\sum_{j} q_{sj}}$
 $t = t + \Delta t$
 $s = k$
record transition from s to k in Δt units

end while