

An Example in SMALL (1/2)

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- Observe that $\mathcal{R}[\text{read}]\rho_0 \sigma_1 = \langle 5, \sigma_1[0/lin] \rangle \equiv \langle 5, \sigma_2 \rangle$.

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- With this, $\mathcal{D}[\text{var } x = \text{read}]\rho_0 \sigma_1 =$
 $(\mathcal{R}[\text{read}]\rho_0 \star \lambda v.\lambda \sigma.\langle \rho_0[\text{new } \sigma/x], \sigma[v/\text{new } \sigma] \rangle)\sigma_1 =$
 $\langle \rho_0[\text{new } \sigma_2/x], \sigma_2[5/\text{new } \sigma_2] \rangle = \langle \rho_0[l/x], \sigma_2[5/l] \rangle \equiv \langle \rho_1, \sigma_3 \rangle$

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- Hence, $\mathcal{C}[c]\rho_0 \sigma_1 = \mathcal{C}[x = x + 1]\rho_1 \sigma_3 =$
 $(\mathcal{E}[x]\rho_1 \star \text{checkLOC} \star \lambda L.\mathcal{R}[x + 1]\rho_1 \star \lambda v.\lambda \sigma.\sigma[v/L])\sigma_3$

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- $\mathcal{R}[[x + 1]]\rho_1 =$
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- $\mathcal{R}[[x]]\rho_1 = \mathcal{E}[[x]]\rho_1 \star \lambda v.\lambda\sigma \mathbf{cases} \dots \mathbf{endcases}$ yields $\mathcal{R}[[x_1]]\rho_1 \sigma_3 =$
 $(\mathbf{cases} \dots \mathbf{endcases})\rho_1(x) \sigma_3 = \langle \sigma_3(\rho_1(x)), \sigma_3 \rangle = \langle 5, \sigma_3 \rangle$

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- Hence, $\mathcal{C}[\![c]\!] \rho_0 \sigma_1 = \mathcal{C}[\![x = x + 1]\!] \rho_1 \sigma_3 =$
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- $\mathcal{R}[\![x + 1]\!] \rho_1 =$
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- $\mathcal{R}[\![1]\!] \rho_1 = \mathcal{E}[\![1]\!] \rho_1 \star \lambda v. \lambda \sigma. \mathbf{cases} \dots \mathbf{endcases}$ and
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- $\mathcal{R}[[1]]\rho_1 = \mathcal{E}[[1]]\rho_1 \star \lambda v. \lambda \sigma. \mathbf{cases} \dots \mathbf{endcases}$ and
 $\mathcal{E}[[1]]\rho_1 = \lambda \sigma. \langle 1, \sigma \rangle$ yield $\mathcal{R}[[1]]\rho_1 \sigma_3 = \langle 1, \sigma_3 \rangle$
- Consequently, $\mathcal{R}[[x + 1]]\rho_1 \sigma_3 = \langle 6, \sigma_3 \rangle$ and we get
 $\lambda v. \lambda \sigma. \sigma[v/l]6 \ \sigma_3 = \sigma_3[6/l]$

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- Consequently, $\mathcal{R}[[x + 1]]\rho_1 \sigma_3 = \langle 6, \sigma_3 \rangle$ and we get
 $\lambda v. \lambda \sigma. \sigma[v/l]6 \ \sigma_3 = \sigma_3[6/l]$
- Therefore, it holds that $\mathcal{C}[[c]]\rho_0 \sigma_1 = \sigma_0[0/lin, nil/lout, 6/l]$ (and
 $\mathcal{P}[[\mathbf{program} \ c]5 :: 0 = nil]$).